Docket: SVL920040039US1 Application: 10/711,808

REMARKS

This is in response to the Final Office Action mailed 2/8/2007, and further in view of the

Request for Continued Examination (RCE) submitted herewith.

Applicants have cancelled previously pending claims 1-22 and 24-25. Claims 25-49 have

been newly added with respect to the RCE concurrently filed herewith. Applicants wish to

emphasize that they are not conceding in this response that those claims are not patentable over

the art cited by the Examiner, as the present claim amendments and cancellations are only for

facilitating expeditious prosecution. Applicants respectfully reserve the right to pursue these and

other claims in one or more continuations and/or divisional patent applications.

This response should obviate outstanding issues and make the pending claims allowable.

Reconsideration of this application is respectfully requested in view of this response.

STATUS OF CLAIMS

Claims 1-22, 24-25 are cancelled.

Claims 25-49 are newly added.

OVERVIEW OF CLAIMED INVENTION

The present invention provides for a computer-based method to <u>version a node rage</u> and

locate a versioned node range in a storage architecture managing node ranges, wherein the

computer-based method comprises the steps of: (a) receiving a node modification request (e.g., a

node deletion request, a node insertion request, a node modification request, etc.) for a node

range from a database system; (b) versioning said node range by copying, to a storage, a node

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range to which said node modification request is to be made and labeling said copied node

range with an identifier; (c) locating said labeled node range via said identifier (e.g., timestamp

or Log Sequence Number) and a hash on said node range; and (d) outputting said located

labeled node range.

The present invention also provides for a computer-based method to version a node

range and to locate a versioned node range in a storage architecture managing node ranges via a

node id range index, said each node assigned a node id value and a set of nodes forming a node

range, each entry in said node id range index pointing to a node range and its range identifier,

RID, said computer-based method implemented in computer readable program code stored in

computer memory, said method comprising the steps of: (a) receiving a node modification

request (e.g., a node deletion request, a node insertion request, a node modification request, etc.)

for a range; (b) versioning said range associated with said node modification request by

shadowing nodes in said range to a Version Hash Table based on RID and assigning a time

identifier to copies of said range; (c) locating a node in said shadowed range via said time

identifier and RIDs; and (d) outputting said located node range.

New readers, after a modification, access current nodes through a new RID and old

readers access old nodes via the same RID, with the shadowed copy being locatable in said

Version Hash Table by hashing the same RID.

In one embodiment, when modifications cause nodes in a range to be moved to a new

RID, previous readers are redirected from the new RID to the old RID via a Redirection Hash

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Table. In another embodiment, when modifications cause nodes in a range to be moved to a new

RID, previous readers are redirected from the new RID to the old RID via an index that describes

where old versions are in said Version Hash Table. In another embodiment, for range deletions,

the range being deleted is moved to reserved RID RIDFF.

The present invention also provides computer medium carrying computer readable

program code implementing the above-described methods.

COMMENTS REGARDING GANESH ET AL. (6,057,236)

AND ODOM ET AL. (6,516,320)

Ganesh teaches a computer-implemented method for providing a data item to a

transaction, wherein the method comprises the steps of: (a) locating, within volatile memory, a

first version of a data block that includes a first version of the data item; (b) determining whether

the first version of the data item is usable by the transaction without respect to whether the first

version of the data block is useable by the transaction; (c) if the first version of the data item is

usable by the transaction, then establishing said data item as a candidate that can be provided to

the transaction; and (d) if the first version of the data item is not useable by the transaction, then

obtaining a version of the data item that is usable by the transaction from a second version of the

data block that is different from said first version.

Odom teaches tiered hashing for data access wherein a memory for access by a program

being executed includes a data access structure stored in memory, the data access structure

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including a first and second index structure together forming a tiered index. According to

Odom, at least one entry in the first structure indicates an entry in the second structure.

Ganesh and Odom fail to disclose, explicitly or implicitly, a storage architecture storing

ranges of nodes. Further absent in either the Ganesh or Odom references, or the combined

teachings of the Ganesh and Odom references, is the versioning of such ranges of nodes.

Also absent in the Ganesh or Odom references, or the combined teachings of the Ganesh

and Odom references, is such versioning of a range of nodes by copying, to a storage, a **node**

range to which said node modification request is to be made and labeling said copied node

range with an identifier, such as a timestamp or log sequence number. Also absent from the

Ganesh and Odom references is a teaching for locating a labeled node range using an

identifier, such as a timestamp or Log Sequence Number, and a hash on said node range;

and outputting such a located labeled node range.

Similarly absent from the Ganesh and Odom references is a teaching for versioning a

range associated with said node modification request by shadowing nodes in said range to a

Version Hash Table based on RID and assigning a time identifier to copies of said range,

wherein the shadowed range is located via said time identifier and RIDs;

Absent such teachings, Ganesh and Odom cannot render obvious Applicants' pending

claims.

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SUMMARY

As has been detailed above, none of the references, cited or applied, provide for the

specific claimed details of Applicants' presently claimed invention, nor renders them obvious. It

is believed that this case is in condition for allowance and reconsideration thereof and early

issuance is respectfully requested.

This response is being filed with a request for extension of time. The Commissioner is

hereby authorized to charge the extension fee, as well as any deficiencies in the fees provided to

Deposit Account No. 09-0460.

If it is felt that an interview would expedite prosecution of this application, please do not

hesitate to contact Applicants' representative at the below number.

Respectfully submitted,

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